

Chapter 4

Finite Difference Method for Parabolic Equations

Last Session Contents:

- 1) Numerical Stability
- 2) Convergence
- 3) Tridiagonal Matrix Algorithm
- 4) Implicit Methods
- 5) Boundary Treatment for Derivative BCs
- 6) Keller-Box Method

1

Numerical Stability

- A concept only defined in iterative problems.
- **It necessitates:**
Errors, of any type, **should not grow** in an iterative process.
- Somewhat more difficult than the study of consistency!
- For non-linear problems, the **necessary condition** for stability is that linear stability analysis of them **must** be stable.
- We will discuss it in detailed later on!
- Now, let's only take a brief look at "stability of **Dufort- Frankel** and **Explicit** scheme"

2

Numerical Stability-In Practice

1. Recall the discretized equation of heat conduction using Dufort-Frankel:

$$\frac{u_i^{n+1} - u_i^{n-1}}{2\Delta t} = \frac{1}{(\Delta x)^2} [u_{i+1}^n - (u_i^{n+1} + u_i^{n-1}) + u_{i-1}^n]$$
 - This scheme is **unconditionally stable**.
2. Explicit Method is stable if:

$$r = \left| \frac{\Delta t}{(\Delta x)^2} \right| \leq \frac{1}{2} \quad \text{It limits time step size!}$$
3. Central Difference in time:

$$\frac{u_i^{n+1} - u_i^{n-1}}{2\Delta t} = \frac{1}{(\Delta x)^2} (u_{i+1}^n - 2u_i^n + u_{i-1}^n)$$
 - This scheme is **Unconditionally Unstable**.

3

Numerical Stability-Physical Interpretation

Sometimes numerical instability can be seen as physically unacceptable results!

Let's consider explicit scheme for discretization of heat equation:

$$u_i^{n+1} = r(u_{i+1}^n + u_{i-1}^n) + (1 - 2r)u_i^n$$

$$r = \frac{\Delta t}{(\Delta x)^2}$$

Assume that at $t = n$ we have: $u_i^n = 0$ and $u_{i+1}^n = u_{i-1}^n = 100^\circ\text{C}$

In this case, if $r > \frac{1}{2}$ temperature at point i will exceed the temperature of two nearby points!

UNACCEPTABLE!?

The maximum expected temperature must be 100°C
However, when $r = 1$ it becomes $T_i^{n+1} = 200^\circ\text{C}$!

4

Convergence

- Generally speaking:
A **Consistent** and **Stable** Scheme will converge!
- Convergence:**
Solving discretized equation of a PDE subjected to similar boundary and initial conditions will converge to the exact solution of that PDE **provided that grid size is chosen to be infinitely small.**
- Finite Difference Equation** is converging if:

$$\lim_{h,k \rightarrow 0} |U^n - u^n| = 0 \quad (x_i, t_n) \in \Omega$$
 where $u^n = u(x_i, t_n)$
- Lax's Equivalence Theorem:**
For a linear well-posed problem, with correct boundary condition, and a Finite Difference Approximation of it, **Consistency and Stability are necessary and sufficient conditions** to provide the convergence!

5

Tridiagonal Systems of Equations

Using numerical methods, the governing PDEs convert to system of algebraic equations as follow:

$$\mathbf{T}\mathbf{x}=\mathbf{b}$$

Large tridiagonal systems arise naturally in a number of problems, especially in the numerical solution of differential equations by **implicit methods.**

$$\mathbf{T} = \begin{bmatrix} a_{11} & a_{12} & 0 & 0 & 0 & \dots & 0 & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 & 0 & \dots & 0 & 0 & 0 \\ 0 & a_{32} & a_{33} & a_{34} & 0 & \dots & 0 & 0 & 0 \\ 0 & 0 & a_{43} & a_{44} & a_{45} & \dots & 0 & 0 & 0 \\ \dots & \dots \\ 0 & 0 & 0 & 0 & 0 & \dots & a_{n-1,n-2} & a_{n-1,n-1} & a_{n-1,n} \\ 0 & 0 & 0 & 0 & 0 & \dots & 0 & a_{n,n-1} & a_{n,n} \end{bmatrix}$$

When a large system of linear algebraic equations has a special pattern, it is usually worthwhile to develop special methods for that unique pattern.

6

Tridiagonal Systems of Equations

One algorithm that deserves special attention is the algorithm for tridiagonal matrices, often referred to as the **Thomas (1949) algorithm.**

$$\mathbf{T} = \begin{bmatrix} a_{11} & a_{12} & 0 & 0 & 0 & \dots & 0 & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 & 0 & \dots & 0 & 0 & 0 \\ 0 & a_{32} & a_{33} & a_{34} & 0 & \dots & 0 & 0 & 0 \\ 0 & 0 & a_{43} & a_{44} & a_{45} & \dots & 0 & 0 & 0 \\ \dots & \dots \\ 0 & 0 & 0 & 0 & 0 & \dots & a_{n-1,n-2} & a_{n-1,n-1} & a_{n-1,n} \\ 0 & 0 & 0 & 0 & 0 & \dots & 0 & a_{n,n-1} & a_{n,n} \end{bmatrix}$$

Row 2: $R_2 - (a_{21}/a_{11})R_1 \rightarrow [0 \quad a_{22} - (a_{21}/a_{11})a_{12} \quad a_{23} \quad 0 \quad 0 \quad \dots \quad 0 \quad 0 \quad 0]$

only a_{32} in column 2 must be eliminated from row 3
 only a_{43} in column 3 must be eliminated from row 4, etc.
 The eliminated element itself does not need to be calculated.
 storing the **elimination multipliers**, $em = (a_{21}/a_{11})$ etc, in place of the eliminated

7

Tridiagonal Systems of Equations

$$\mathbf{T} = \begin{bmatrix} a_{11} & a_{12} & 0 & 0 & 0 & \dots & 0 & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 & 0 & \dots & 0 & 0 & 0 \\ 0 & a_{32} & a_{33} & a_{34} & 0 & \dots & 0 & 0 & 0 \\ 0 & 0 & a_{43} & a_{44} & a_{45} & \dots & 0 & 0 & 0 \\ \dots & \dots \\ 0 & 0 & 0 & 0 & 0 & \dots & a_{n-1,n-2} & a_{n-1,n-1} & a_{n-1,n} \\ 0 & 0 & 0 & 0 & 0 & \dots & 0 & a_{n,n-1} & a_{n,n} \end{bmatrix}$$

Hint:
 Only the diagonal element in each row is affected by the elimination.
 Elimination in rows 2 to n is accomplished as follows:

$$a_{i,i} = a_{i,i} - (a_{i,i-1}/a_{i-1,i-1})a_{i-1,i} \quad (i = 2, \dots, n)$$

Thus, the elimination step involves only **2n** multiplicative operations to place T in upper triangular form.

8

Tridiagonal Systems of Equations

Subsequent elements of the **b** vector are changed in a similar manner.

$$b_2 = b_2 - (a_{21}/a_{11})b_1$$

a_{21} is already calculated. Thus, the total process of elimination, including the operation on the **b** vector, requires only **3n** multiplicative operations.

Hint:
The $n \times n$ tridiagonal matrix **T** can be stored as an $n \times 3$ matrix **A'** since there is no need to store the zeros.

$$A' = \begin{bmatrix} - & a'_{1,2} & a'_{1,3} \\ a'_{2,1} & a'_{2,2} & a'_{2,3} \\ a'_{3,1} & a'_{3,2} & a'_{3,3} \\ \dots & \dots & \dots \\ a'_{n-1,1} & a'_{n-1,2} & a'_{n-1,3} \\ a'_{n,1} & a'_{n,2} & - \end{bmatrix}$$

Column1=Sub-diagonal elements of T
Column2=Diagonal elements of T
Column3=Super-diagonal elements of T

9

Example

$$T'' - \alpha^2 T = -\alpha^2 T_a \quad T(0.0) = 0.0 \quad T(1.0) = 100.0$$

$$\frac{T_{i+1} - 2T_i + T_{i-1}}{\Delta x^2} + 0(\Delta x^2) - \alpha^2 T_i = -\alpha^2 T_a$$

$$T_{i-1} - (2 + \alpha^2 \Delta x^2)T_i + T_{i+1} = -\alpha^2 \Delta x^2 T_a$$

$$\begin{cases} \alpha = 4.0 \\ \Delta x = 0.125 \\ (2 + \alpha^2 \Delta x^2) = 2.25 \end{cases}$$

$$A' = \begin{bmatrix} - & -2.25 & 1.0 \\ 1.0 & -2.25 & 1.0 \\ 1.0 & -2.25 & 1.0 \\ 1.0 & -2.25 & 1.0 \\ 1.0 & -2.25 & 1.0 \\ 1.0 & -2.25 & 1.0 \\ 1.0 & -2.25 & - \end{bmatrix} \quad \text{and} \quad b = \begin{bmatrix} 0.0 \\ 0.0 \\ 0.0 \\ 0.0 \\ 0.0 \\ 0.0 \\ -100.0 \end{bmatrix}$$

10

Example

$$A' = \begin{bmatrix} - & -2.250000 & 1.0 \\ (-0.444444) & -1.805556 & 1.0 \\ (-0.553846) & -1.696154 & 1.0 \\ (-0.589569) & -1.660431 & 1.0 \\ (-0.602253) & -1.647747 & 1.0 \\ (-0.606889) & -1.643111 & 1.0 \\ (-0.608602) & -1.641398 & - \end{bmatrix} \quad \text{and} \quad b' = \begin{bmatrix} 0.0 \\ 0.0 \\ 0.0 \\ 0.0 \\ 0.0 \\ 0.0 \\ -100.0 \end{bmatrix}$$

$$x_7 = b_7/a'_{7,2} = (-100)/(-1.641398) = 60.923667$$

$$x_6 = (b_6 - a'_{6,3}x_7)/a'_{6,2} = [0 - (1.0)(60.923667)]/(-1.643111) = 37.078251$$

$$x = \begin{bmatrix} 1.966751 \\ 4.425190 \\ 7.989926 \\ 13.552144 \\ 22.502398 \\ 37.078251 \\ 60.923667 \end{bmatrix}$$

11

Implicit Finite Difference Approximations

- Backward Difference Scheme**

$$\frac{u_i^{n+1} - u_i^n}{\Delta t} = \frac{u_{i-1}^{n+1} - 2u_i^{n+1} + u_{i+1}^{n+1}}{\Delta x^2} + O(\Delta t, \Delta x^2)$$

Considering $r = \frac{k}{h^2}$ we have:

$$u_i^{n+1} - u_i^n = r(u_{i-1}^{n+1} - 2u_i^{n+1} + u_{i+1}^{n+1})$$

or,

BTCS: $-ru_{i-1}^{n+1} + (1 + 2r)u_i^{n+1} - ru_{i+1}^{n+1} = u_i^n$

How to solve it?!

12

Implicit Finite Difference Approximations

BTCS: $-ru_i^{n+1} + (1 + 2r)u_i^{n+1} - ru_{i+1}^{n+1} = u_i^n$

- Assume that boundary values are zero at both ends.
- This tri-diagonal system can be solved by Thomas Algorithm.

Note that:

- BTCS is **unconditionally stable**.
- Second order in space but first order in time!

$$\begin{bmatrix} (1+2r) & -r & & & \\ -r & (1+2r) & -r & & \\ & \ddots & \ddots & \ddots & \\ & & -r & (1+2r) & -r \\ & & & -r & (1+2r) \end{bmatrix} \begin{bmatrix} u_2^{n+1} \\ u_3^{n+1} \\ \vdots \\ \vdots \\ u_{max-1}^{n+1} \end{bmatrix} = \begin{bmatrix} u_2^n \\ u_3^n \\ \vdots \\ \vdots \\ u_{max-1}^n \end{bmatrix}$$

Implicit Finite Difference Approximations

- Crank-Nicolson Scheme**

$$\frac{1}{\Delta t} \delta_t u_i^{n+\frac{1}{2}} = \frac{1}{2(\Delta x)^2} (\delta_x^2 u_i^{n+1} + \delta_x^2 u_i^n)$$

$$\frac{u_i^{n+1} - u_i^n}{\Delta t} = \frac{1}{2(\Delta x)^2} \times (u_{i-1}^{n+1} - 2u_i^{n+1} + u_{i+1}^{n+1} + u_{i-1}^n - 2u_i^n + u_{i+1}^n)$$

Considering $r = \frac{k}{\alpha \Delta t}$ we have:

$$u_i^{n+1} - u_i^n = \frac{r}{2}(u_{i-1}^{n+1} - 2u_i^{n+1} + u_{i+1}^{n+1}) + \frac{r}{2}(u_{i-1}^n - 2u_i^n + u_{i+1}^n)$$

or,

$$ru_{i-1}^{n+1} - 2(1+r)u_i^{n+1} + ru_{i+1}^{n+1} = -ru_{i-1}^n + 2(-1+r)u_i^n - ru_{i+1}^n$$

Crank-Nicolson

Implicit Finite Difference Approximations

- Keller-Box Scheme**

$$\frac{\partial u}{\partial t} = \frac{\partial^2 u}{\partial x^2}$$

We can re-write this equation as:

$$\frac{\partial u}{\partial x} = P$$

$$\frac{\partial u}{\partial t} = \frac{\partial P}{\partial x}$$

Grid Stencil

Implicit Finite Difference Approximations

- Keller-Box Scheme**

$$\frac{\partial u}{\partial t} = \frac{\partial^2 u}{\partial x^2}$$

We can re-write this equation as:

$$\frac{\partial u}{\partial x} = P$$

$$\frac{\partial u}{\partial t} = \frac{\partial P}{\partial x} \rightarrow \frac{u_i^n - u_{i-1}^n}{h_i} = P_{i-\frac{1}{2}}^n$$

or,

$$u_i^n - u_{i-1}^n - \frac{h_i}{2}(P_i^n + P_{i-1}^n) = 0$$

Grid Stencil

Implicit Finite Difference Approximations

- Keller-Box Scheme**

$$\frac{\partial u}{\partial t} = \frac{\partial^2 u}{\partial x^2}$$

We can re-write this equation as:

$$\frac{\partial u}{\partial x} = p$$

$$\frac{\partial u}{\partial t} = \frac{\partial p}{\partial x} \rightarrow \frac{p_i^{n+1} - p_i^n}{h_t} = \frac{u_{i+1}^{n+1} - u_{i-1}^{n+1}}{k_x}$$

Grid Stencil

× Unknowns
 □ Knowns
 ○ "Centering"

$$\frac{1}{2} \left[\frac{p_i^n - p_{i-1}^n}{h_t} + \frac{p_i^{n+1} - p_{i-1}^{n+1}}{h_t} \right] = \frac{1}{2} \left[\frac{u_i^n + u_{i-1}^n}{k_x} - \frac{u_i^{n+1} + u_{i-1}^{n+1}}{k_x} \right]$$

or,

$$P_i^n - P_{i-1}^n - \frac{h_t}{k_x} (u_i^n + u_{i-1}^n) = P_{i-1}^{n+1} - P_i^{n+1} - \frac{h_t}{k_x} (u_i^{n+1} + u_{i-1}^{n+1})$$

$$\underbrace{\hspace{10em}}_{R_{i-\frac{1}{2}}^{n-\frac{1}{2}}}$$

Implicit Finite Difference Approximations

$$\begin{cases} u_i^n - u_{i-1}^n - \frac{h_t}{2} (P_i^n + P_{i-1}^n) = 0 \\ P_i^n - P_{i-1}^n - r_i^n (u_i^n + u_{i-1}^n) = R_{i-\frac{1}{2}}^{n-\frac{1}{2}} \end{cases}$$

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ -1 & -\frac{h_t}{2} & 1 & -\frac{h_t}{2} \\ -r_i^n & -1 & -r_{i-1}^n & 1 \\ 0 & 0 & -1 & -\frac{h_t}{2} & 1 & -\frac{h_t}{2} \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \end{bmatrix} \begin{bmatrix} u_0 \\ P_0 \\ u_1 \\ P_1 \\ u_2 \\ P_2 \\ \vdots \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ R_1 \\ 0 \\ R_2 \\ 0 \\ \vdots \end{bmatrix}$$

Implicit Finite Difference Approximations

- We can re-write the previous matrix as below where the elements are blocks itself.
- This matrix can be solved using **block Thomas algorithm**.
- Please note that this matrix should be constructed so that: $\det(B_0) \neq 0$

$$\begin{bmatrix} B_0 & C_0 & & & & \\ A_1 & B_1 & C_1 & & & \\ & A_2 & B_2 & C_2 & & \\ & & \ddots & \ddots & \ddots & \\ & & & A_n & B_n & \\ & & & & & \end{bmatrix} \begin{bmatrix} u_0 \\ u_1 \\ u_2 \\ \vdots \\ u_n \end{bmatrix} = \begin{bmatrix} R_0 \\ R_1 \\ R_2 \\ \vdots \\ R_n \end{bmatrix}$$

Implicit Finite Difference Approximations

- The main features of Keller Box Scheme

- Only slightly more arithmetic to solve than the Crank-Nikolson method
- Second order accurate with arbitrary (uniform) x and y spacing
- Allows very rapid x variation
- Allows easy programming of the solution of large numbers of coupled equation.

- Steps:
 - Reduce the Equations to a 1st - order system
 - Write difference equations using central differencing.
 - Linearize the resulting algebraic equation and write them in matrix-vector form
 - Solve the linear system by the block-tridiagonal elimination method

Implementation of Boundary Condition

- Implicit schemes mostly end up to this form:

$$A_i u_{i-1}^{n+1} + B_i u_i^{n+1} + C_i u_{i+1}^{n+1} = R_i^n = R_i(u_{i-1}^n, u_i^n, u_{i+1}^n)$$

$$u_1 = u_a, \quad u_{imax} = u_b$$
- 1st Method:**
Boundary conditions are considered in the equations and matrix-form equation is solved for $i = 2$ to $i = imax - 1$.

$$A_2 u_1^{n+1} + B_2 u_2^{n+1} + C_2 u_3^{n+1} = R_2^n \quad i = 2$$

$$B_2 u_2^{n+1} + C_2 u_3^{n+1} = R_2^n - A_2 u_1^{n+1} \quad \text{Considering BC at } a$$

$$A_{imax-1} u_{imax-2}^{n+1} + B_{imax-1} u_{imax-1}^{n+1} + C_{imax-1} u_{imax}^{n+1} = R_{imax-1}^n \quad i = imax$$

$$A_{imax-1} u_{imax-2}^{n+1} + B_{imax-1} u_{imax-1}^{n+1} = R_{imax-1}^n - C_{imax-1} u_{imax}^{n+1} \quad \text{Considering BC at } imax$$

21

Implementation of Boundary Condition

- 1st Method:**
Boundary conditions are considered in the equations and matrix-form equation is solved for $i = 2$ to $i = imax - 1$.

$$B_2 u_2^{n+1} + C_2 u_3^{n+1} = R_2^n - A_2 u_1^{n+1}$$

$$A_{imax-1} u_{imax-2}^{n+1} + B_{imax-1} u_{imax-1}^{n+1} = R_{imax-1}^n - C_{imax-1} u_{imax}^{n+1}$$

$$\begin{bmatrix} B_2 & C_2 \\ A_1 & B_1 & C_1 \\ \vdots & \vdots & \vdots \\ A_{i-2} & B_{i-2} & C_{i-2} \\ A_{i-1} & B_{i-1} & \end{bmatrix} \begin{bmatrix} u_2^{n+1} \\ u_3^{n+1} \\ \vdots \\ u_{i-2}^{n+1} \\ u_{i-1}^{n+1} \end{bmatrix} = \begin{bmatrix} R_2^n - A_2 u_1^{n+1} \\ R_1^n \\ \vdots \\ R_{i-2}^n \\ R_{i-1}^n - C_{i-1} u_{imax}^{n+1} \end{bmatrix}$$

22

Implementation of Boundary Condition

- 2nd Method:**
Let the computer do the calculations!

$$\begin{bmatrix} B_1 & C_1 \\ A_2 & B_2 & C_2 \\ \vdots & \vdots & \vdots \\ A_{i-1} & B_{i-1} & C_{i-1} \\ A_i & B_i & \end{bmatrix} \begin{bmatrix} u_1^{n+1} \\ u_2^{n+1} \\ \vdots \\ u_{i-1}^{n+1} \\ u_i^{n+1} \end{bmatrix} = \begin{bmatrix} R_1^n \\ R_2^n \\ \vdots \\ R_{i-1}^n \\ R_i^n \end{bmatrix}$$

$$B_1 = B_i = 1 \quad C_1 = A_i = 0$$

Note: a slight increase in computation cost, however, gives more flexibility in computer code!

23

Derivative Boundary Condition

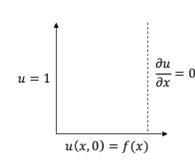
- 1st Method:**
Backward difference at the boundary

$$\frac{\partial u}{\partial x} \Big|_{i=N} = \frac{3u_N - 4u_{N-1} + u_{N-2}}{2\Delta x} = 0$$

$$u_N = \frac{4u_{N-1} - u_{N-2}}{3}$$
- 2nd Method:**
False boundary

$$\frac{\partial u}{\partial x} \Big|_{i=N} = \frac{u_{N+1} - u_{N-1}}{2\Delta x} = 0$$

$$u_{N+1} = u_{N-1}$$



24

Numerical Solution of Blasius Equation

- Blasius Equation:**

$$f'''(\eta) + \frac{\alpha}{2} f(\eta) f''(\eta) = 0 \quad \eta = y \sqrt{\frac{U_\infty}{\alpha \nu x}} \quad \alpha \text{ an arbitrary parameter}$$

$$f(0) = f'(0) = 0, \quad f'(\infty) = 1$$
- Breaking it up to three first order equations:**

$$\frac{df}{d\eta} = u \quad f(0) = 0$$

$$\frac{du}{d\eta} = v \quad u(0) = 0 \quad u(\infty) = 1$$

$$\frac{dv}{d\eta} = -\frac{\alpha}{2} f v$$

25

Numerical Solution of Blasius Equation

- We discretize the equations in $\eta_{j-\frac{1}{2}}$**

$$\frac{f_j - f_{j-1}}{h_j} = u_{j-1/2} = \frac{1}{2}(u_j + u_{j-1})$$

$$\frac{u_j - u_{j-1}}{h_j} = v_{j-1/2} = \frac{1}{2}(v_j + v_{j-1})$$

$$\frac{v_j - v_{j-1}}{h_j} = -\frac{\alpha}{2} (f v)_{j-1/2} = -\alpha \frac{f_{j-1/2} v_{j-1/2}}{4}$$
- Newton Linearization**
 These equations are **non-linear**, so, we have to linearize them.

$$f_k^{n+1} = f_k^n + \delta f_k^n$$

$$u_k^{n+1} = u_k^n + \delta u_k^n$$

$$v_k^{n+1} = v_k^n + \delta v_k^n$$
 where n denotes the iteration number.
Note: we call the solution converged if $\delta(\cdot)$ variables approach to zero!

26

Numerical Solution of Blasius Equation

- Substituting these parameters into the first equations yields:**

$$f_j^n + \delta f_j^n - f_{j-1}^n - \delta f_{j-1}^n = \frac{h_j}{2} [u_j^n + \delta u_j^n + u_{j-1}^n + \delta u_{j-1}^n]$$

$$u_j^n + \delta u_j^n - u_{j-1}^n - \delta u_{j-1}^n = \frac{h_j}{2} [v_j^n + \delta v_j^n + v_{j-1}^n + \delta v_{j-1}^n]$$

$$v_j^n + \delta v_j^n - v_{j-1}^n - \delta v_{j-1}^n = -\frac{\alpha h_j}{4} [(f_j^n + \delta f_j^n)(v_j^n + \delta v_j^n) + (f_{j-1}^n + \delta f_{j-1}^n)(v_{j-1}^n + \delta v_{j-1}^n)]$$
- We can rewrite it as:**

$$\delta f_j^n - \delta f_{j-1}^n - \frac{h_j}{2} (\delta u_j^n + \delta u_{j-1}^n) = r_j^n \quad r_j^n = f_{j-1}^n - f_j^n + h_j u_{j-1/2}^n$$

$$\delta u_j^n - \delta u_{j-1}^n - \frac{h_j}{2} (\delta v_j^n + \delta v_{j-1}^n) = s_j^n \quad s_j^n = u_{j-1}^n - u_j^n + h_j v_{j-1/2}^n$$

$$(1 + \frac{\alpha h_j}{4} f_j^n) \delta v_j^n + (-1 + \frac{\alpha h_j}{4} f_{j-1}^n) \delta v_{j-1}^n = t_j^n \quad t_j^n = v_{j-1}^n - v_j^n - \frac{\alpha h_j}{2} (f v)_{j-1/2}^n$$

27

Numerical Solution of Blasius Equation

- Finally, it can be written in matrix form as:**

$$A_j^n \delta_{j-1}^n + B_j^n \delta_j^n + C_j^n \delta_{j+1}^n = \vec{R}_j^n \quad \delta_j^n = [\delta f_j^n \quad \delta u_j^n \quad \delta v_j^n]^T$$

$$\begin{bmatrix} B_1^n & C_1^n \\ A_2^n & B_2^n & C_2^n \\ & \ddots & \ddots & \ddots \\ & & A_{j_{max}-1}^n & B_{j_{max}-1}^n & C_{j_{max}-1}^n \\ & & & A_{j_{max}}^n & B_{j_{max}}^n & C_{j_{max}}^n \end{bmatrix} \begin{bmatrix} \delta_1^n \\ \delta_2^n \\ \vdots \\ \delta_{j_{max}-1}^n \\ \delta_{j_{max}}^n \end{bmatrix} = \begin{bmatrix} \vec{R}_1^n \\ \vec{R}_2^n \\ \vdots \\ \vec{R}_{j_{max}-1}^n \\ \vec{R}_{j_{max}}^n \end{bmatrix}$$
- Note:** This block-tridiagonal matrix can be solved using block Thomas elimination

28





Numerical Solution of Blasius Equation

- *A, B, C* and *R* blocks are as following:

$$A_j^* = \begin{bmatrix} -1 & -\frac{h_j}{2} & 0 \\ \frac{\alpha h_j}{4} s_{j-1}^* & 0 & -1 + \frac{\alpha h_j}{4} f_{j-1}^* \\ 0 & 0 & 0 \end{bmatrix} \quad 2 \leq j \leq j_{max}$$

$$B_j^* = \begin{bmatrix} 1 & -\frac{h_j}{2} & 0 \\ \frac{\alpha h_j}{4} s_j^* & 0 & 1 + \frac{\alpha h_j}{4} f_j^* \\ 0 & -1 & -\frac{h_{j+1}}{2} \end{bmatrix} \quad 2 \leq j \leq j_{max} - 1$$

$$C_j^* = \begin{bmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 1 & -\frac{h_{j+1}}{2} \end{bmatrix} \quad 1 \leq j \leq j_{max} - 1$$

$$R_j^* = \begin{bmatrix} f_j^* \\ s_j^* \\ f_{j+1}^* \end{bmatrix} \quad 2 \leq j \leq j_{max} - 1$$

- Boundary condition implementation also gives:

$$B_1^* = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -1 & -\frac{h_2}{2} \end{bmatrix}$$

$$B_{j_{max}}^* = \begin{bmatrix} 1 & -\frac{h_{j_{max}}}{2} & 0 \\ \frac{\alpha h_{j_{max}}}{4} s_{j_{max}}^* & 0 & 1 + \frac{\alpha h_{j_{max}}}{4} f_{j_{max}}^* \\ 0 & 1 & 0 \end{bmatrix}$$

$$R_1^* = \begin{bmatrix} 0 \\ 0 \\ f_2^* \end{bmatrix}$$

$$R_{j_{max}}^* = \begin{bmatrix} f_{j_{max}}^* \\ s_{j_{max}}^* \\ 0 \end{bmatrix}$$

29